

1. The application of abstract data types:
 - (a) Demands much discipline even in a programming language such as Java
 - (b) Is possible only when a programming language such as Java is used
 - (c) Happens automatically when you use the programming language Java

2. The relation between *abstract data types* and (*Java*) *interfaces* is:
 - (a) They are two different names for the same concept
 - (b) The concept 'interface' is an extension of the concept 'abstract data type'
 - (c) Interfaces are suitable to express abstract data types in Java

3. During the determination of the 'Big-Oh' notation for the *running time* $T(n)$ of an algorithm, we can neglect the actual time spent by an elementary operation, because:
 - (a) We are only interested in an approximation
 - (b) It does not influence the asymptotic behavior of the algorithm
 - (c) The actual time spent cannot be measured

4. Which running time indicates a *polynomial* algorithm?
 - (a) $T(n) = O(n \log(n))$
 - (b) $T(n) = O(n^2 + n^3)$
 - (c) $T(n) = O(2^n)$
 - (d) $T(n) = O(\log(n))$

5. What is the asymptotic *worst-case* running time (in Big-Oh notation) of the following variant of Bubblesort algorithm:

```

void sort(String A[], int n) {
    boolean bubbled = false;
    for (int i=1; i<n; i++) {
        bubbled = false;
        for (int j=0; j<n-i; j++) {
            if (A[j] >= A[j+1]) {
                String y = A[j]; A[j] = A[j+1]; A[j+1] = y;
                bubbled = true;
            }
        }
        if (!bubbled) break;
    }
}

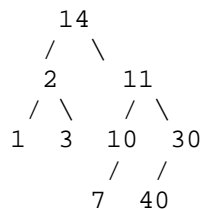
```

- (a) $T(n) = O(n)$
 - (b) $T(n) = O(2n + n \log(n))$
 - (c) $T(n) = O(n^2)$
 - (d) $T(n) = O(n \log(n))$
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6. What is the running time of the algorithm in question 5 in the case that array $A[]$ is already sorted (ascending)?
 - (a) $T(n) = O(n)$
 - (b) $T(n) = O(2n + n \log(n))$
 - (c) $T(n) = O(n^2)$
 - (d) $T(n) = O(n \log(n))$

 7. Assume, a *Stack* ADT is (efficiently) implemented with a single Queue. What can you tell about the complexity of $Push()$ and $Pop()$?
 - (a) Both operations have to be $O(n)$ (n is the number of elements on the Stack)
 - (b) One of them can be $O(1)$ and the other one has to be $O(n)$
 - (c) One of them can be $O(n)$ and the other one has to be $O(n^2)$

8. Observe the concepts *position* and *rank of an element* in a Sequence based on an Array:
- Both the *rank* and the *position* of an element agree with the index of the element in the Array
 - The *rank* of an element agrees with the index of the element in the Array, the *position* is relative and is determined with the aid of the index.
 - The *position* of an element agrees with the index of the element in the Array, the *rank* is relative and is determined with the aid of the index.
 - Neither the *position* nor the *rank* have any relation with the index of the Array

9. What is the *in-order* list of the elements in the following binary tree?



- 14 – 2 – 1 – 3 – 11 – 10 – 7 – 30 – 40
- 1 – 3 – 2 – 7 – 10 – 40 – 30 – 11 – 14
- 1 – 2 – 3 – 14 – 7 – 10 – 11 – 40 – 30
- 1 – 2 – 3 – 7 – 10 – 11 – 14 – 30 – 40

10. We want to use an Euler-tour to compute an expression (e.g.: $3*5 / 2 + 5$) that is represented in a binary tree. The Euler-tour is executed by the *Template-method pattern*. In which of the following four methods do we perform the operations (+, -, /, *, etc.)?

- VisitLeft
- VisitBelow
- VisitRight
- VisitExternal

11. The *Heap-order property* means that:

- The left child of a node always has a larger value than the right child
- All children at the left side of a node have smaller values than the children at the right side of a node
- Children do not have smaller values than their parents
- Nodes at depth d always have larger values than nodes at depth $d-1$

12. During the insertion of an element in a Binary Search Tree,

- The element is always placed in an external node that becomes internal
- First, room is made in an internal node by moving some elements
- One or more *restructures* are performed after the element is placed

13. AVL trees and Splay-trees both make use of rotations (*tri-node restructure operations*). In AVL trees, this can happen at insertions and deletions, in Splay trees it happens more than once at *all operations*.

- It means that Splay trees are always less preferable than AVL trees because AVL trees need less rotations
- The height of AVL trees is guaranteed to be $O(\log n)$; with Splay trees this is not the case. So, AVL trees are to be preferred.
- Splay trees are preferable if some elements are used more often than others
- Splay trees are better balanced than AVL trees and thus are preferable.

14. To sort small numbers of items:
- (a) You have to use a quadratic algorithm
 - (b) It is unwise to use an $O(n \log n)$ sorting algorithm
 - (c) You can better use Insertion-Sort than Quick-Sort
 - (d) It does not matter which algorithm is used
15. The cut-off value in Quick-Sort indicates:
- (a) The spot where the list to sort is split in two
 - (b) The maximal size of a list to sort on which Quick-Sort still is used
 - (c) The minimal size of a list to sort on which Quick-Sort still is used
 - (d) The maximal depth of the recursion
16. Dynamic programming is to be preferred over Divide-and-Conquer when:
- (a) There are many subproblems
 - (b) Subproblems do overlap much
 - (c) The problem cannot be divided into subproblems
 - (d) The solution of a problem is not easily obtained from solutions of subproblems
17. The advantage of the Hirschberg algorithm for finding the *Optimal Alignment* of two DNA-strings over the *Longest Common Substring*-method from the handbook is:
- (a) Hirschberg applies Divide-and-Conquer and on dynamic programming
 - (b) The running time of the Hirschberg-algorithm is much lower
 - (c) (Much) less storage space is needed for the Hirschberg-algorithm
 - (d) The algorithms are in fact identical. There is no advantage
18. A *back edge* appears during depth-first search in a directed graph if:
- (a) A node does not have children
 - (b) A child of a node already has been visited and also is a descendent of that node
 - (c) A child of a node already has been visited but is not a descendent of that node
 - (d) An arrow crosses an arrow that is already visited before
19. Finding substring T of string S in a *Suffix Trie* of S costs (not including the construction of the Suffix Trie):
- (a) $O(|S|)$ operations
 - (b) $O(|T|)$ operations
 - (c) $O(|S|+|T|)$ operations
 - (d) $O(|S| \times |T|)$ operations
20. In a *region-based k-D tree*:
- (a) Areas are divided in two equal parts on each internal node
 - (b) The distribution of points in the area determines the way in which the area is split
 - (c) The size of the area determines whether an area is split

Open question:

Describe three different applications of *Tries* in Knowledge Engineering.